

Communication Complexity of Gap Hamming Problem à la Sherstov

Board talk based on

<http://eccc.hpi-web.de/report/2011/063/>

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The Gap Hamming Problem

Alice: $x \in \{0, 1\}^n$

Bob: $y \in \{0, 1\}^n$

$$\text{Answer} = \begin{cases} \text{Yes} & \text{if } d(x, y) > \frac{n}{2} + \sqrt{n} \\ \text{No} & \text{if } d(x, y) < \frac{n}{2} - \sqrt{n} \\ \text{Don't care} & \text{Otherwise} \end{cases}$$

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More convenient: $x, y \in \{-1, +1\}^n$.

$$\text{Answer} = \begin{cases} -1 & \text{if } \langle x, y \rangle \leq -\sqrt{n} \\ +1 & \text{if } \langle x, y \rangle \geq \sqrt{n} \\ \text{Don't care} & \text{Otherwise.} \end{cases}$$

Gap Orthogonality (Sherstov)

Alice: $x \in \{-1, +1\}^n$.

Bob: $y \in \{-1, +1\}^n$.

$$\text{Answer} = \begin{cases} -1 & \text{if } |\langle x, y \rangle| \leq \sqrt{n}/8 \\ +1 & \text{if } |\langle x, y \rangle| \geq \sqrt{n}/4 \\ \text{Don't care} & \text{Otherwise} \end{cases}$$

Concentration

Theorem (Talagrand)

Let V be a linear subspace of \mathbb{R}^n . There is a constant $c > 1$ such that for a random $x \in \{-1, +1\}^n$ and all $t > 0$,

$$\Pr[|\|\pi_V x\| - \sqrt{\dim V}| > t + c] < 4 \exp\left(-\frac{t^2}{c}\right).$$

Almost orthogonal vectors

Lemma

Fix $A \subseteq \{-1, +1\}^n$ sufficiently large ($|A| > 2^{(1-\alpha)n}$). Then, one can find $m > k/10$ elements $x_1, x_2, \dots, x_m \in A$ such that

$$\|\pi_{S_{i-1}}(x_i)\| \leq \frac{\sqrt{n}}{3}.$$

Projection on to almost orthogonal vectors

Lemma

Fix vectors x_1, x_2, \dots, x_m are almost orthogonal. Then,

$$\Pr_{y \in \{-1, +1\}^n} \left[\forall i \left| \langle y, x_i \rangle \right| \leq \frac{\sqrt{n}}{4} \right] \leq \exp(-\Omega(m)).$$

No large red rectangle

Theorem

Let $\gamma > 0$ be small enough. Let $R = A \times B$ be a rectangle.

Suppose

$$\Pr_{(x,y) \in R} \left[|\langle x, y \rangle| > \frac{\sqrt{n}}{4} \right] \leq \gamma.$$

Then,

$$\mu(R) = 4^{-n} |A| |B| < \exp(-\Omega(n)).$$